

Approximate Inference by Semidefinite Relaxations

Andrea Montanari

[with Adel Javanmard, Federico Ricci-Tersenghi, Subhabrata Sen]

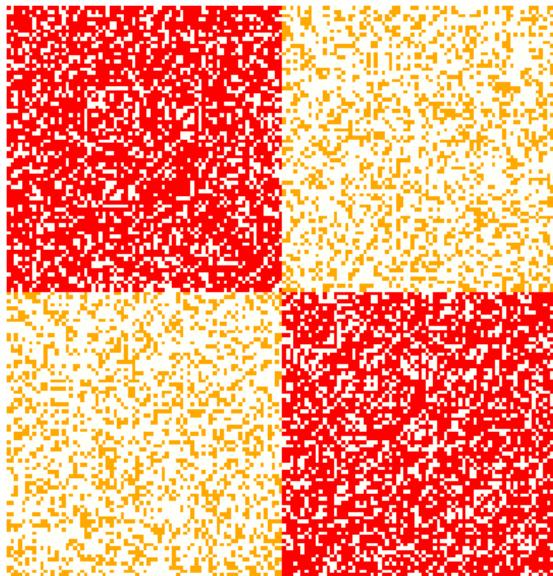
Stanford University

December 11, 2015

What is this talk about?

SDP for Matrix/Graph estimation

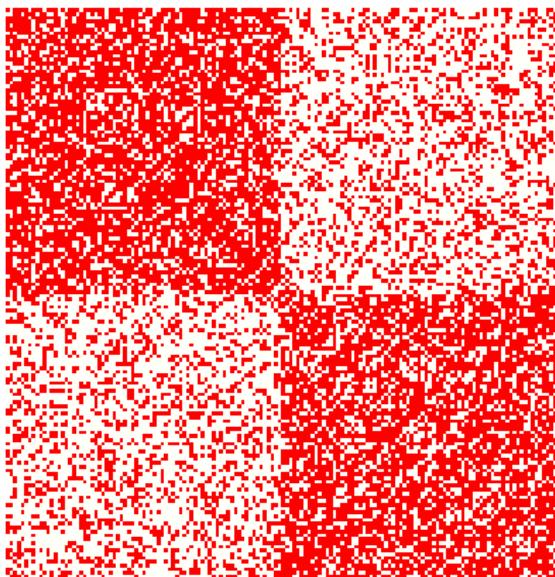
The hidden partition model



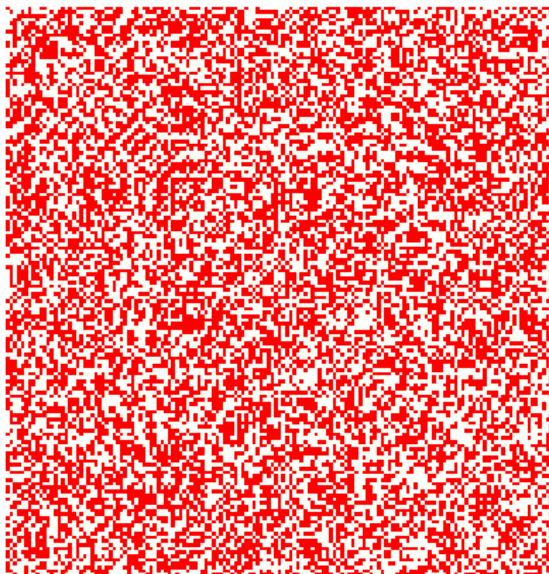
Vertices V , $|V| = n$, $V = V_+ \cup V_-$, $|V_+| = |V_-| = n/2$

$$\mathbb{P}\{(i,j) \in E\} = \begin{cases} p & \text{if } \{i,j\} \subseteq V_+ \text{ or } \{i,j\} \subseteq V_-, \\ q < p & \text{otherwise.} \end{cases}$$

Of course entries are not colored...



... and rows/columns are not ordered



Problem: Detect/estimate the partition

What is this talk about?

SDP for Matrix/Graph estimation

Exact phase transition(?)

Outline

1 Background

2 Near-optimality of SDP

3 How does SDP work ‘in practice’?

4 Conclusion

Background

Statistical estimation

$$x_{0,i} = \begin{cases} +1 & \text{if } i \in V_+, \\ -1 & \text{if } i \in V_-, \end{cases}$$

$$\mathbb{P}\{(i,j) \in E\} = \begin{cases} p & \text{if } x_{0,i} = x_{0,j}, \\ q < p & \text{otherwise.} \end{cases}$$

Estimator $\hat{\mathbf{x}} \in \{+1, -1\}^n$

$$\text{Overlap}_n(\hat{\mathbf{x}}) = \frac{1}{n} \mathbb{E}\{|\langle \hat{\mathbf{x}}(G), \mathbf{x}_0 \rangle|\}.$$

Statistical estimation ($p = a/n$, $q = b/n$)

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Information theory threshold

Theorem (Mossel, Neeman, Sly, 2012)

There is an estimator that achieves $\liminf_{n \rightarrow \infty} \text{Overlap}_n(\hat{\mathbf{x}}) \geq \varepsilon > 0$ if and only if $a + b > 2$ and

$$\frac{a - b}{\sqrt{2(a + b)}} > 1.$$

[Proves conjecture by Decelle, Krzakala, Moore, Zdeborova, 2011]

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Computational threshold

- ▶ Dyer, Frieze 1989 $p = na > q = nb$ fixed.
- ▶ Condon, Karp 2001 $a - b \gg n^{1/2}$
- ▶ McSherry 2001 $a - b \gg \sqrt{b \log n}$
- ▶ Coja-Oghlan 2010 $a - b \gg \sqrt{b}$
- ▶ Massoulie 2013 and Mossel, Neeman, Sly, 2013

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Very ingenious spectral methods!

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What if I am not ingenious?

Maximum Likelihood

Posterior probability

Candidate partition $\sigma \in \{+1, -1\}^n$

$$\mathbb{P}(x_0 = \sigma | G) \approx \frac{1}{Z(G)} \prod_{(i,j) \in E} \{a \mathbb{I}(\sigma_i = \sigma_j) + b \mathbb{I}(\sigma_i \neq \sigma_j)\} \mathbb{I}\left(\sum_{i=1}^n \sigma_i = 0\right)$$

Pairwise binary graphical model

Adjacency matrix

$$A_{ij} = \begin{cases} 1 & \text{if } (i, j) \in E, \\ 0 & \text{otherwise.} \end{cases}$$

$$\mathbf{A} = (A_{ij})_{1 \leq i, j \leq n}$$

Maximum likelihood

$$\sigma_i = \begin{cases} +1 & \text{if } i \in V_+, \\ -1 & \text{if } i \in V_-. \end{cases}$$

$$\text{maximize} \quad \sum_{i,j=1}^n A_{ij} \sigma_i \sigma_j,$$

$$\text{subject to} \quad \sum_{i=1}^n \sigma_i = 0,$$

$$\sigma_i \in \{+1, -1\}.$$

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Lagrangian

$$\begin{aligned} & \text{maximize} && \sum_{i,j=1}^n A_{ij} \sigma_i \sigma_j - \gamma \left(\sum_{i=1}^n \sigma_i \right)^2. \\ & \text{subject to} && \sigma_i \in \{+1, -1\}. \end{aligned}$$

A good choice:

$$\gamma = \frac{a+b}{2n} \equiv \frac{d}{n}$$

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Centered adjacency matrix

$$A_{ij}^{\text{cen}} = \begin{cases} 1 - (d/n) & \text{if } (i, j) \in E, \\ -(d/n) & \text{otherwise.} \end{cases}$$

$$\mathbf{A}^{\text{cen}} = \mathbf{A} - \frac{d}{n} \mathbf{1} \mathbf{1}^T$$

Lagrangian

$$\begin{aligned} & \text{maximize} && \langle A^{\text{cen}}, \sigma \sigma^T \rangle, \\ & \text{subject to} && \sigma \in \{+1, -1\}^n. \end{aligned}$$

- ▶ NP-hard
- ▶ $\text{SDP}(A^{\text{cen}})$ is a very natural convex relaxation

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Relaxation

$$\begin{aligned} & \text{maximize} && \langle A^{\text{cen}}, \sigma\sigma^T \rangle, \\ & \text{subject to} && \sigma \in \{+1, -1\}^n. \end{aligned}$$

SDP(A^{cen}):

$$\begin{aligned} & \text{maximize} && \langle A^{\text{cen}}, X \rangle, \\ & \text{subject to} && X \in \mathbb{R}^{n \times n}, \quad X \succeq 0, \\ & && X_{ii} = 1. \end{aligned}$$

Estimator

- ▶ Compute principal eigenvector $v_1(X)$
- ▶ Threshold it $\hat{x}^{\text{SDP}}(G) = \text{sign}(v_1(X))$
- ▶ Randomized variation for proofs

This is really off-the-shelf

How well does it work?

Estimator

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How well does it work?

Near-optimality of SDP

Before we pass to SDP

- ▶ What's the problem with sparse graphs?
- ▶ What's the problem vanilla PCA?

Why PCA?

Ground truth

$$x_{0,i} = \begin{cases} +1 & \text{if } i \in V_+, \\ -1 & \text{if } i \in V_-. \end{cases}$$

Data = RankOne + Wigner

$$\frac{1}{\sqrt{d}} A^{\text{cen}} = \frac{\lambda}{n} x_0 x_0^\top + W, \quad \lambda \equiv \frac{a - b}{\sqrt{2(a + b)}}$$

$$E\{W_{ij}\} = 0, \quad \mathbb{E}\{W_{ij}^2\} \in \left\{ \frac{a}{dn}, \frac{b}{dn} \right\} \approx \frac{1}{n}.$$

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The *right* parametrization

$$d = \frac{a + b}{2}, \quad \lambda = \frac{a - b}{\sqrt{2(a + b)}}$$

Naive PCA

$$\hat{\mathbf{x}}^{\text{PCA}}(A^{\text{cen}}) = \sqrt{n} v_1(A^{\text{cen}}).$$

Does it work?

$$\frac{1}{\sqrt{d}} \mathbf{A}^{\text{cen}} = \frac{\lambda}{n} \mathbf{x}_0 \mathbf{x}_0^\top + \mathbf{W}$$

Naive idea:

$$\|\mathbf{W}\|_2 \leq \text{const.}, \quad \left\| \frac{\lambda}{n} \mathbf{x}_0 \mathbf{x}_0^\top \right\|_2 = \lambda \Rightarrow \text{Works for } \lambda = O(1)$$

Wrong!

Does it work?

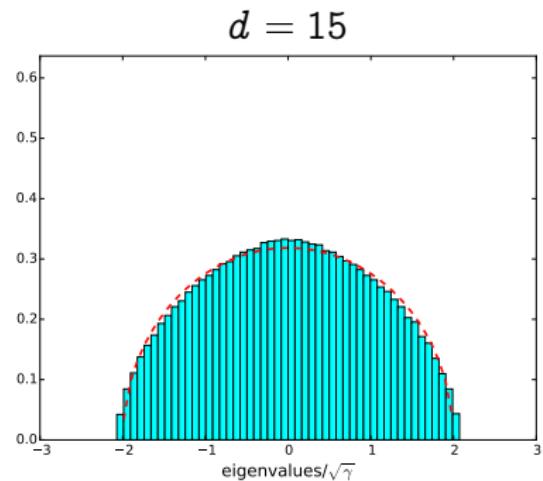
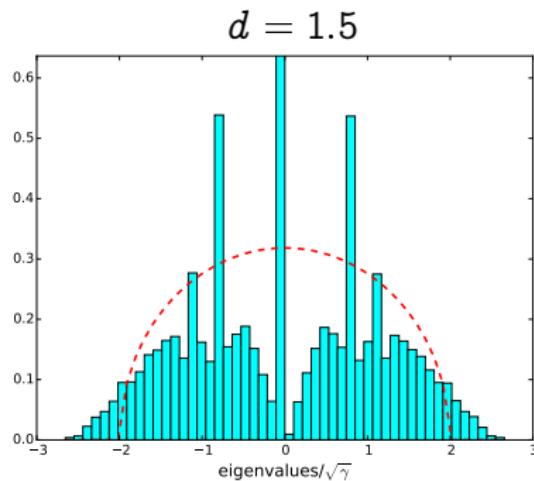
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Spectral relaxation bad in the sparse regime!

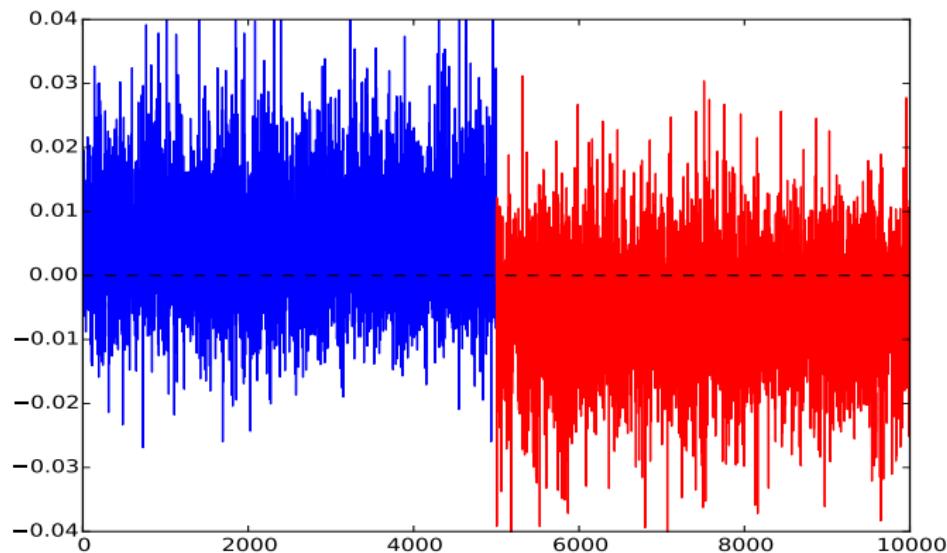


Theorem (Krivelevich, Sudakov 2003+Vu 2005)

With high probability,

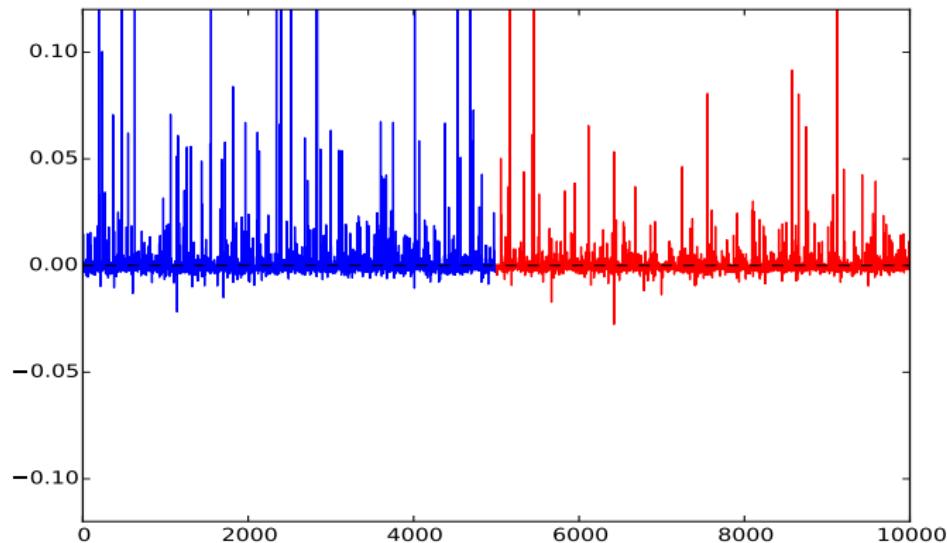
$$\lambda_{\max}(A^{\text{cen}}/\sqrt{d}) = \begin{cases} 2(1 + o(1)) & \text{if } d \gg (\log n)^4, \\ C \sqrt{\log n / (\log \log n)}(1 + o(1)) & \text{if } d = O(1). \end{cases}$$

Example: $d = 20$, $\lambda = 1.2$, $n = 10^4$



$$v_1(A^{\text{cen}})$$

Example: $d = 3$, $\lambda = 1.2$, $n = 10^4$



$$v_1(A^{\text{cen}})$$

Why should SDP work better?

$$\begin{aligned} & \text{maximize} && \langle \mathbf{A}^{\text{cen}}, \mathbf{X} \rangle, \\ & \text{subject to} && \mathbf{X} \in \mathbb{R}^{n \times n}, \mathbf{X} \succeq 0, \\ & && \mathbf{X}_{ii} = 1. \end{aligned}$$

Recall the ultimate limit

$G(n, d, \lambda)$ graph distribution with parameters

$$d = \frac{a + b}{2} > 1, \quad \lambda = \frac{a - b}{\sqrt{2(a + b)}}$$

Theorem (Mossel, Neeman, Sly, 2012)

If $\lambda < 1$, then

$$\lim \sup_{n \rightarrow \infty} \|G(n, d, 0) - G(n, d, \lambda)\|_{\text{TV}} < 1.$$

If $\lambda > 1$, then

$$\lim_{n \rightarrow \infty} \|G(n, d, 0) - G(n, d, \lambda)\|_{\text{TV}} = 1.$$

SDP has nearly optimal threshold

Theorem (Montanari, Sen 2015)

Assume $G \sim G(n, d, \lambda)$.

If $\lambda \leq 1$, then, with high probability,

$$\frac{1}{n\sqrt{d}} \text{SDP}(A_G^{\text{cen}}) = 2 + o_d(1).$$

If $\lambda > 1$, then there exists $\Delta(\lambda) > 0$ such that, with high probability,

$$\frac{1}{n\sqrt{d}} \text{SDP}(A_G^{\text{cen}}) = 2 + \Delta(\lambda) + o_d(1).$$

Consequence

Corollary (Montanari, Sen 2015)

Assume $\lambda \geq 1 + \varepsilon$. Then there exists $d_0(\varepsilon)$ and $\delta(\varepsilon) > 0$ such that the randomized SDP-based estimator achieves, for $d \geq d_0(\varepsilon)$,

$$\liminf_{n \rightarrow \infty} \mathbb{E}\{\text{Overlap}_n(\hat{x}^{\text{SDP}})\} \geq \delta(\varepsilon).$$

Earlier/related work

Optimal spectral tests

- ▶ Massoulie 2013
- ▶ Mossel, Neeman, Sly, 2013
- ▶ Bordenave, Lelarge, Massoulie, 2015

SDP, $d = \Theta(\log n)$

- ▶ Abbe, Bandeira, Hall 2014
- ▶ Hajek, Wu, Xu 2015

SDP, detection

- ▶ Guédon, Vershynin, 2015 (requires $\lambda \geq 10^4$, very different proof)

How does SDP work ‘in practice’?

Thresholds

- ▶ $\lambda_c^{\text{opt}}(d) \equiv$ Threshold for optimal test
- ▶ $\lambda_c^{\text{SDP}}(d) \equiv$ Threshold for SDP-based test

What we know

- ▶ $\lambda_c^{\text{opt}}(d) = 1$ [Mossel, Neeman, Sly, 2013]
- ▶ $\lambda_c^{\text{SDP}}(d) = 1 + o_d(1)$ [Montanari, Sen, 2015]

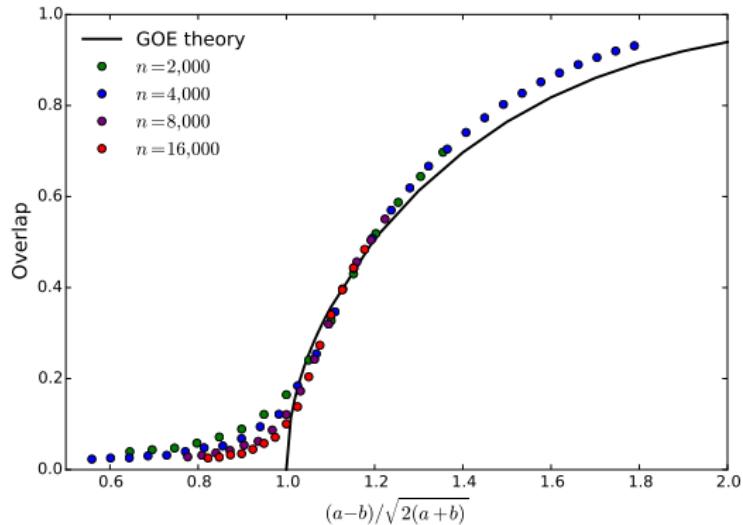
How big is the $o_d(1)$ gap?

What we know

- ▶ $\lambda_c^{\text{opt}}(d) = 1$ [Mossel, Neeman, Sly, 2013]
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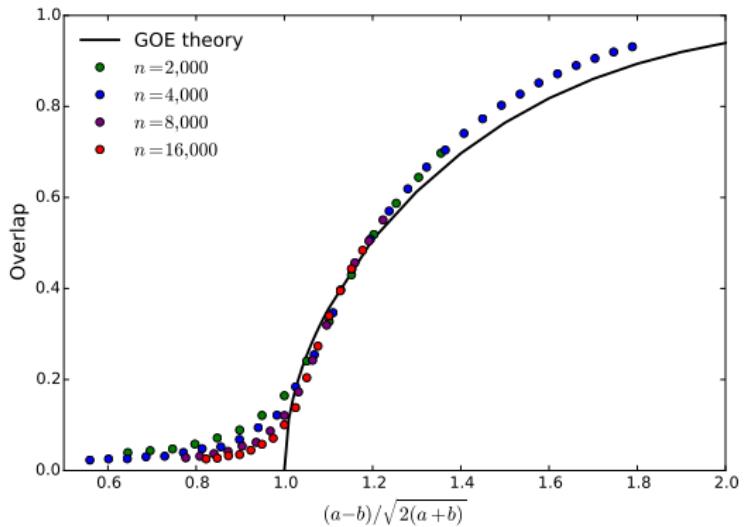
Simulations: $d = 5$, $N_{\text{sample}} = 500$ (with Javanmard and Ricci)



SDP estimator $\hat{x}^{\text{SDP}} \in \{+1, -1\}^n$

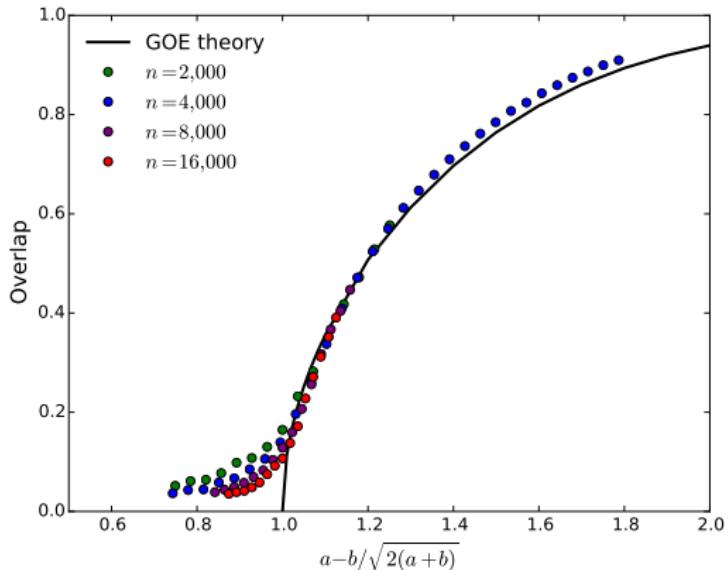
$$\text{Overlap}_n(\hat{x}) = \frac{1}{n} \mathbb{E}\{|\langle \hat{x}^{\text{SDP}}(G), x_0 \rangle|\}.$$

Simulations: $d = 5$, $N_{\text{sample}} = 500$



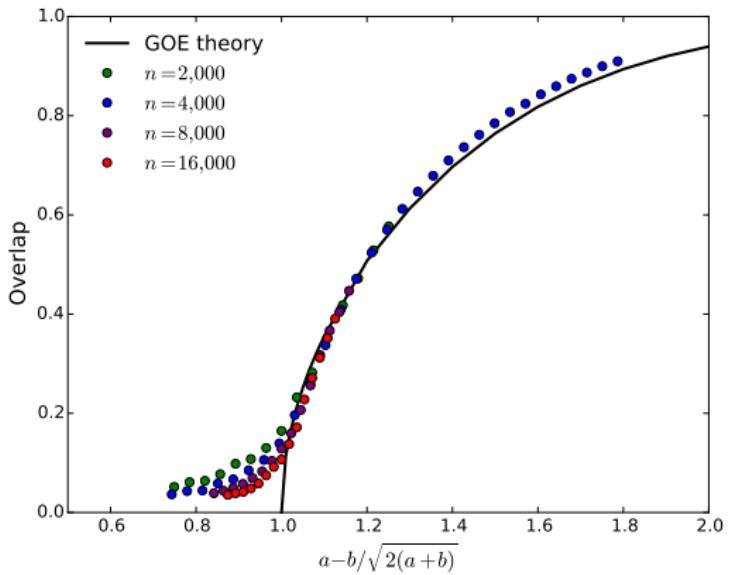
$$\lambda_c^{\text{SDP}}(d = 5) \approx 1.$$

Simulations: $d = 10$, $N_{\text{sample}} = 500$



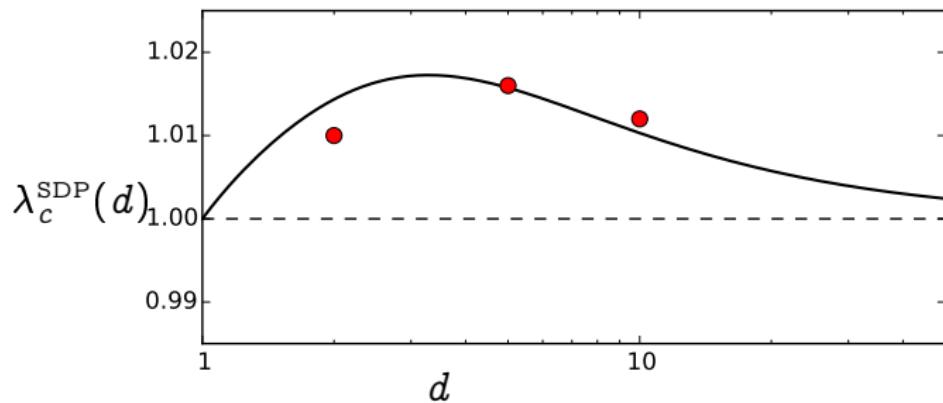
$$\lambda_c^{\text{SDP}}(d = 10) \approx 1 .$$

Simulations: $d = 10$, $N_{\text{sample}} = 500$



Can we estimate $\lambda_c^{\text{SDP}}(d)$ from data?

$$\lambda_c^{\text{SDP}}(d), N_{\text{sample}} \geq 10^5 \quad (\text{10 years CPU time})$$



- ▶ Dots: Numerical estimates
- ▶ Line: Non-rigorous analytical approximation
(using statistical physics)
- ▶ **At most 2% sub-optimal!**

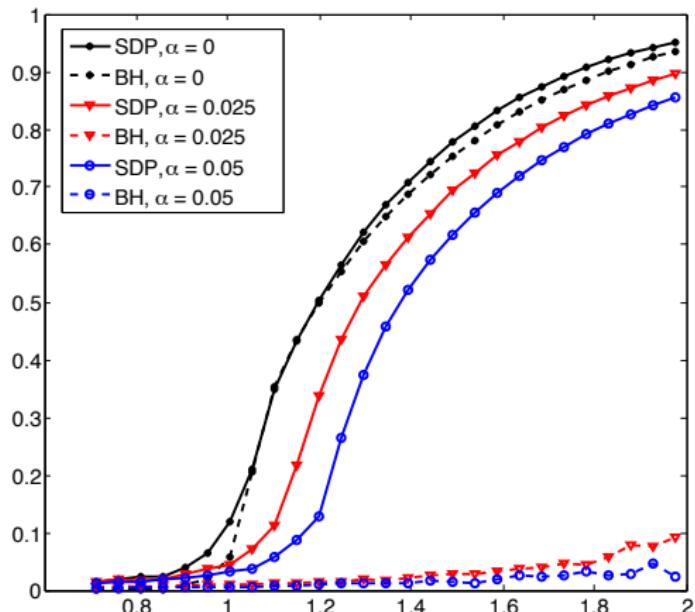
One last question

Is this approach robust to model miss-specifications?

An experiment

- ▶ Select $S \subseteq V$ uniformly at random. with $|S| = n\alpha$.
- ▶ For each $i \in S$, connect all of its neighbors.

An experiment



- ▶ Solid line: SDP
- ▶ Dashed line: Spectral
(Non-backtracking walk [Krzakala, Moore, Mossel, Neeman, Sly, Zdeborova, Zhang, 2013])

Conclusion

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- ▶ SDP \gg PCA when data are heterogeneous
- ▶ Sharp information about eigenvalues of random matrices
- ▶ A lot of work on SDP with random data
 - [Srebro, Fazel, Parrillo, Candés, Recht, Gross, myself, ...]
- ▶ Little known about ‘sharp SDP properties’ and SDP vs PCA

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